

Finding Primes

$$\pi(N) \sim \frac{N}{\log N}.$$

of primes $< N$

for integers of size $\sim N$, probability of $\frac{1}{\log N}$ of being prime.

Finding primes reduces to primality testing.

Primality Testing

Fermat's Little Theorem: Let p prime. Let a be coprime to p .

Then $a^{p-1} \equiv 1 \pmod{p}$.

Fermat Primality Test

Let $n > 1$ be a number to test.

Choose $1 < a < n$. (Could test $\gcd(a, n)$.)

Compute $a^{n-1} \pmod{n}$.

If $a^{n-1} \not\equiv 1$ then n is not prime.

If $a^{n-1} \equiv 1$ then n is "probably prime".

Ex. ① $n = 1007$

Choose $a = 2$

$$\text{Compute } 2^{1006} \pmod{1007} \\ \equiv 271$$

$\Rightarrow$ not prime

② $n = 1009$

Choose $a = 2$

$$\text{Compute } 2^{1008} \pmod{1009} \\ \equiv 1$$

$\Rightarrow$ probably prime

Try $a = 3$

$$3^{1008} \pmod{1009} \\ \equiv 1$$

$\Rightarrow$ very probably prime.

③ $n = 561$

Choose $a = 2$

$$\text{Compute } 2^{560} \pmod{561} \\ \equiv 1$$

$\Rightarrow$ probably prime

ACK! $3 \mid 561$.

$$3^{560} \not\equiv 1 \quad (\text{since } 3 \text{ not div.})$$

Notes: ① For more confidence, try more a 's.

② Stronger versions, e.g. Miller-Rabin, Solovay-Strassen.

How many "fakers" (false positives).

For context, $X < 10^{10}$

3 455 052 511 primes

14 884

pseudoprimes to base $a=2$

(pass Fermat test w/ $a=2$).

3 291

strong pseudoprimes to base $a=2$

(pass Miller-Rabin w/ $a=2$)

0 strong pseudoprimes to base 2, 3, 5, 7
simultaneously.

A Carmichael number is an integer n s.t. for all $0 < a < n$,
if $\gcd(a, n) = 1$ then $a^{n-1} \equiv 1 \pmod{n}$ but n is composite.

We know a lower bd for # of Carmichael #'s up to X of $X^{1/3}$.

In practice we use "probably prime" (probabilistic) tests
without any issue in cryptographic ranges.

Square roots & factoring

Sps $x^2 \equiv y^2 \pmod{n}$ but $x \not\equiv \pm y \pmod{n}$

Then $(x+y)(x-y) \equiv x^2 - y^2 \equiv 0 \pmod{n}$
 $\not\equiv 0 \quad \not\equiv 0$

↙ not 1 or n

So $x+y$ and $x-y$ each contain proper factors of n
but not n itself.

So $\gcd(x+y, n)$ and $\gcd(x-y, n)$ should be proper factors of n .

$\left(\begin{array}{c} \text{finding interesting} \\ \text{square roots} \\ \text{mod } n \end{array} \right) \implies \left(\begin{array}{c} \text{factoring} \\ n \end{array} \right)$

Ex. $n = 295927$.

If we notice $544^2 \equiv 3^2 \pmod{n}$.

$$\gcd(n, 544 - 3) = 541$$

$$\gcd(n, 544 + 3) = 547$$

In fact, $n = 541 \cdot 547$.

Fermat factoring: If $n = x^2 - y^2$ then $n = (x+y)(x-y)$.

try $n + 1^2 \stackrel{?}{=} \square$

$$n + 2^2 \stackrel{?}{=} \square$$

$$n + 3^2 \stackrel{?}{=} \square \quad (544^2 \text{ in example}) \rightarrow \text{go to example.}$$

This method detects $n = pq$ where p, q are very close in size

Moral: don't choose p, q similar in RSA.

Quadratic Sieve (to factor n)

Factor base of small primes: $2, 3, 5, \dots, p_k < B$.

Obtain relations:

take $k \sim \sqrt{n}$

try to factor (in factor base):

$$\begin{aligned} k^2 - n \\ (k+1)^2 - n \\ (k+2)^2 - n \\ \text{etc.} \end{aligned}$$

Ex.

$$n = 2201$$

$$47^2 - n = 2^3$$

$$49^2 - n = 2^3 \cdot 5^2 \Rightarrow (47 \cdot 49)^2 \equiv 2^3 \cdot 2^3 \cdot 5^2$$

$$\equiv 2^6 \cdot 5^2 \pmod{n}$$

$$\begin{array}{r} (3, 0) \\ + (3, 2) \\ \hline (6, 2) \end{array} \pmod{\varphi(n)}$$